

# Preventing Clone Attacks Using Dynamic Cryptography in MANETs

S. Anusha<sup>1</sup>, Jazira Banu<sup>2</sup>, Mamatha<sup>3\*</sup>, Nikhitha<sup>4</sup>, P. C. Rashmi<sup>5</sup>

<sup>1,2,3,4</sup>Student, Dept. of Computer Science and Engineering, Srinivas School of Engineering, Mangalore, India <sup>5</sup>Assistant Professor, Dept. of Computer Science and Engg., Srinivas School of Engineering, Mangalore, India \*Corresponding author: mamthasherigar@gmail.com

Abstract: Mobile Ad-Hoc Networks (MANETs) use anonymous routing protocols that hide node identifies and/or routes from outside observers in order to provide anonymity protection. In this attack, an adversary captures few nodes, replicates them and then deploys arbitrary number of replicas throughout the network. It is very hard to distinguish between non-compromised nodes a clone node since a clone has the same security and code information of original node. Hence cloned nodes can launch a variety of other attacks. The detection of cloning attacks is therefore a fundamental problem. The main objective of this method is to prevent clone attacks in MANETs while routing. A clone attack is launched by a malicious node by copying the configuration of a legitimate active node of current operating network. Polynomial bivariate keys are used for preventing clone attacks while routing by validating the next hop until the data reaches the destination. A polynomial key based authentication layer in AODV is used to demonstrate the clone attack prevention in the MANET.

Keywords: Clone attack, MANET.

## 1. Introduction

Mobile Ad-hoc Networks (MANETs) fall under the most widely used communication technologies especially because of the features like mobility, dynamic infrastructure, dynamic and easy link establishment and reconfiguration [1]. The mobile devices are free to move randomly and arrange themselves randomly. The communication takes place in MANET by using multi-hop paths. Nodes in the MANET share the wireless medium and the topology of the network changes dynamically. Since the nodes are mobile, there is a chance for communication link breakage. The application areas of MANET have mostly been in critical areas where the routing performance and security is expected to be close to ideal. However, this is not achievable practically and hence strategies to achieve greater security adopt.

The primary concerned security issue in MANETs is to protect the network layer from malicious attacks and to detect and prevent malicious nodes in the communication network [2]. The security solution is important to protect both route and data forwarding operations in the network layer. Lacking any proper security solution, each malicious node tends to act as a readily available router, which will solely disturb the network operation from correct delivering of the packets and the malicious nodes can give stale routing updates or drop all the packets passing through them.

One important mechanism to detect clone attacks is the time domain detection. Time is divided into equal length intervals and the time intervals are associated with the challenge. Trusted node broadcasts the challenge to every node in the network initially. Based on the one-way property of hash function, it can easily verify the authenticity by using any of the previously verified challenge of the preloaded one. Also space domain detection is used for detecting node replication attacks. This scheme consists of two phases: the local check phase and the local witness check phase. The local check is the phase when two needs meet each other and exchange information according to the local information exchange. The witness nodes record the inevitable information during exchanging the information in the node to node. Once the nodes meet each other, they exchange the recorded information about identity.

Clone attack or node replication attack is a severe attack in MANET [3]. In this attack an adversary captures only a small number of nodes replicates them and then deploys arbitrary number of replicas throughout the network. It is very hard to distinguish between non-compromised nodes a clone node since a clone has the same security and code information of original node. Hence cloned nodes can launch a variety of other attacks. Detection and prevention of cloning attacks in a mobile ad hoc network is a fundamental problem and cannot be easily handled [4]. Most of the existing protocols expose the following limitations: high performance overheads, necessity of central control, unreasonable assumptions, lack of smart attack detection etc.

The polynomial key verification is used to differentiate the nodes within the network whether they are replicas or legitimate nodes. On encountering malicious nodes, data is sent through alternate routes to reach the destination. The nodes in the network operate with greater security and hence there is greater packet reception lesser loss in the network. Also by routing through legitimate nodes the energy drain is reduced in the network the reason being there is greater energy remaining in PK-AODV.



#### 2. Related Work

SAR (Secure-Aware Ad Hoc Routing protocol), using AODV, it proposed encryption and decryption process using a common key [5]. SAR defines a level of trust as a metric for routing and as an attribute for security for routing. The main drawback with SAR protocol is whenever the levels of security rise; it needs different keys for different levels, thereby increasing the number of keys [6]. SEAD is a secure efficient ad hoc distance vector routing protocol. It uses one-way hash function without the usage of asymmetric cryptographic mechanism. Authentication is used by this mechanism to differentiate between malicious and non-malicious nodes, and reduces resource consumption attacks launched by malicious nodes. It can overcome DoS, many types of routing attacks and resource consumption attacks also avoid routing loops. The drawback lies whenever the attacker uses the same metric and sequence number used for authentication were same by the recent update message and updates with new update message. ARAN (Authenticated Routing for Ad Hoc Networks) based on cryptographic certificates which overcome all types of attacks in the network layer [7]. ARAN provides authentication, integrity and non-repudiation. However, this protocol mechanism is quite robust against attacks, and is mainly based on prior security coordination among nodes which cannot be correctly assures at all times [8]. Drawbacks for any protocol to be used in the WSN-resource-constrained environment also vulnerability is an issue. Randomized, Efficient, and Distributed (RED) protocol proposed detection of node replication attacks, and self-healing. Detecting replication attacks is a nontrivial problem in MANETs due to the challenges resulted from node mobility, cloned/compromised node conspiracy, and the large number and wide spread replicas [9]. Existing approaches either fails in mobile environments due to the limitations caused by local views or their dependence on invariant claims such as location and neighbor list, or are constrained by the number and malicious activities of the replicas. Detecting Replica Attacks in Mobile Ad Hoc Networks proposed two replication detection schemes (TDD and SDD) to tackle all these challenges from both the time domain and the space domain [10]. The authors have proved that TDD and SDD provide high detection accuracy and excellent resilience against smart and colluding replicas, have no restriction on the number and distribution of replicas, and incur low communication/computation overhead. Dynamic detection of node replication attack aims to detect the cloned node in the environment network. The most obvious attack in wireless sensor network is node replication attack [11]. This attack the nodes are replicated manually based on their id and key values. Cloned node or adversary promotes the node key or id of the legitimate node, creates more replicas of the particular node in the current network with the same id and also this node may cripple the entire network. It detects clone replication attack in dynamic way and to detect the replicas in mobile wireless sensor network [12]. Existing schemes rely on fixed

sensor locations and hence do not work in mobile sensor networks, where sensors are expected to move [13]. A fast and effective mobile replica node detection scheme using the sequential probability ratio test was developed [14]. The problem of replica node attacks in MANETs is tackled here. Using SPRT, this scheme detects mobile replicas in an efficient and robust manner at the cost of reasonable overheads. Executing and checking the test cases is not possible all the time in the manual testing. Therefore, selecting a test case and ranking is important [15]. The objective of the test case selection is to have a better test case from a pool of test cases and assigning the rank to each test case will leads the software as an error free and which gives a good efficiency. Ranking of test cases is especially useful if a system is a having large number of test cases. Hence selecting and ranking a test case plays an important role in the software testing. Automation tools helps in design and execution of test scripts saving time and cost involved in manual testing [16].

#### 3. Proposed Method

As the topology of the network changes dynamically in the MANET, the key is also calculated dynamically.

The same key is used for the data transmission from X to Y and from Y to X. Even though the attacker compromises the node and capture the key, then introduce the malicious node with same configuration by using the compromised parameters, the rest of the normal nodes present in the network can easily identify it is a replicated node since the key of the changed dynamically.

## A. Flat topology

The flowchart and algorithm description of PK-AODV is shown in figure below.

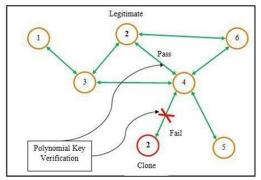


Fig. 1. Polynomial key verification between nodes

Figure 1 shows the operation of the nodes inside the network. All nodes are verified using the polynomial key verification process. If the nodes pass the verification, then the nodes can continue communicating with the verified node. Otherwise, although the node id remains the same, through the polynomial key verification between the nodes 2 and 4, the node 2 is identified as clone of the legitimate node 2. Preventing the cloned attacker, the communication can be performed through



an alternate neighbor of the source node. This process is continuously and consecutively performed until the data reaches the destination.

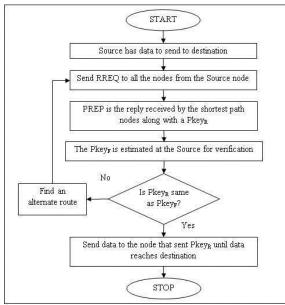


Fig. 2. Flowchart

# B. Hierarchical topology

The flowchart and algorithm description under hierarchical topology is described below.

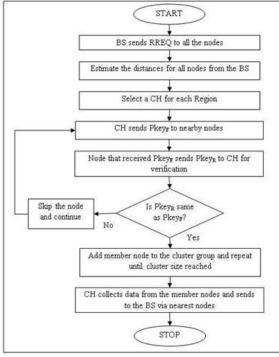


Fig. 3. Flowchart

If the node that detects the malicious node broadcasts information by sending malicious node ID, then the legitimate

node that has the same ID may be eliminated from the routing tables of the neighboring nodes. This is a threat to the network as some of the links will be deleted although they seem to be functioning well. Hence all communicating nodes are verified hop-by-hop whether they are legitimate. This method can hence act as an efficient method to perform communication between the nodes in the network preventing the occurrence of clone attacks.

	Algorithm
Set this_r	node = source;
{	
Get neigh	hbours(this_node);
Send req	uest RREQ to all the neighbours(this_node);
Eachnoo	de sends PREP along with the Pkevg:
If { Pkey Send dat Set this_t	eceives Pkey <sub>R</sub> and estimates Pkey <sub>F</sub> ; y <sub>R</sub> == Pkey <sub>F</sub> } { a this_node → next_node node = next_node; p 1 until data reaches destination;
	node and goto step 4;
	node and goto step 4; Algorithm
Skip the } }	
Skip the } Line No	Algorithm
Skip the } Line No 1	Algorithm For all n e N {
Skip the } } Line No 1 2 3 4	Algorithm For all n e N { BS sends Pilot signal to n;
Skip the } Line No 1 2 3	Algorithm For all n e N { BS sends Pilot signal to n;
Skip the } } Line No 1 2 3 4 5 6	Algorithm For all n e N { BS sends Pilot signal to n; n replies to BS with distances; }
Skip the: } } Line No 1 2 3 4 5	Algorithm         For all n e N {         BS sends Pilot signal to n;         n replies to BS with distances;         )         Pick a node for each region r e G as CH {
Skip the } } Line No 1 2 3 4 5 6	Algorithm         For all n e N {         BS sends Pilot signal to n;         n replies to BS with distances;         }         Pick a node for each region r e G as CH {         If i's distance < Range(CH) { // i e N
Skip the } Line No 1 2 3 4 5 6 7	Algorithm         For all n e N {         BS sends Pilot signal to n;         n replies to BS with distances;         )         Pick a node for each region r e G as CH {         If i's distance < Range(CH) { // i e N
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Skip the: } Line No 1 2 3 4 5 6 7 8 9	Algorithm         For all n e N {         BS sends Pilot signal to n;         n replies to BS with distances;         )         Pick a node for each region r e G as CH {         If i's distance < Range(CH) { // i e N
Skip the: } Line No 1 2 3 4 5 6 7 8 9 10 11 12	Algorithm           For all $n \in \mathbb{N}$ {           BS sends Pilot signal to $n$ ; $n$ replies to BS with distances; $\}$ Pick a node for each region $r \in G$ as CH {           If i's distance < Range(CH) { // $i \in \mathbb{N}$ CH send Pkeys to node i $node$ i replies with Pkeys,           If { Pkeys, = Pkeys} {           Add node i to cluster $r$ until max cluster size is reached
Skip the: } Line No 1 2 3 4 5 6 7 8 9 10 11	Algorithm         For all n e N {         BS sends Pilot signal to n; n replies to BS with distances; }         Pick a node for each region r e G as CH {         If i's distance < Range(CH) { // i e N

Clone Server 1	
Clone Server 1	
Status Configuration	
Node Status	
Node 1 is InActive Node 2 is InActive	Node Status
Node 3 is InActive	
Node 4 is InActive	
Node Comminication Session Information	
	ê
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Select A Node Select	
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a t librar a contractor	
ne Based Approach Node 1	
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Clone Base Browse Reset all Tost Statuton Nodes Hode Registration tect a Destination : Select Total	d Approach Node 1 Transmit Data BBEQ Request Coher Text Node Dynamic Value : 0 1 2 3 4 5 6 7 8 9 10
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itatus				
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Clone Based Approach 1 File Browse Plain Yest Destination Hodes Select a Destination :- Shortest Path :-	Node Registartion	Browse Reset	Transmit Data Cipher Text	RREQ Request

### 4. Conclusion

The polynomial key verification is hence used to differentiate the nodes within the network whether they are replicas or legitimate nodes while communication is performed. On encountering malicious nodes, data is sent through alternate routes to reach the destination. The nodes in the network operate with greater security and hence there is greater packet reception lesser loss in the network. Also by routing through legitimate nodes the energy drain is reduced in the network which is why there is greater energy remaining in PK-AODV.

Future works aim at providing data integrity by combining PK- AODV with security algorithms. Also this mechanism can be adopted by hybrid networks by implementing the suitable

integration modules, which can be done in future work

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