

# Wireless Ad-hoc Networks Performance Evaluation Using NS-2 and NS-3 Network Simulators

J. Sateesh<sup>1</sup>, M. Harika<sup>2</sup>, P. Sirish Kumar<sup>3</sup>

<sup>1</sup>Assistant Professor, Department of Electronics and Communication Engineering, University College of Engineering, JNTUK, Vizianagaram, India

<sup>2</sup>Assistant Professor, Department of Electronics and Communication Engineering, Sarada Institute of Science, Technology and Management, Srikakulam, India

<sup>3</sup>Assistant Professor, Department of Electronics and Communication Engineering, Aditya Institute of Technology and Management, Tekkali, India

*Abstract*: An Ad-hoc network is a collection of wireless terminals that are able to dynamically form a temporary network without any aid from fixed infrastructure or centralized administration. In recent years, another class of adhoc networks are continuing to attract the attention for their potential use in several fields. Mobility and the absence of any fixed infrastructure make mobile Ad-hoc network very attractive for mobility and rescue operations and time-critical applications. In this paper, we concentrate on performance evaluation of wireless ad-hoc network using ns-2 and ns-3 network simulators. We compared the throughput simulation results of ns-2 and ns-3 with experimental results for Linear Topology (LT). We also compared by simulations the performance of ns-2 and ns-3 considering memory size metric.

*Keywords*: Ad-hoc networks; ns-3; ns-2; Linear Topology, Throughput.

#### 1. Introduction

Wireless ad hoc networks are collections of wireless nodes, that communicate directly over a common wireless channel. The nodes are equipped with wireless transceiver. They don't need any additional infrastructure, such as base station or wired access point, etc. Therefore, each node doesn't only plays the role of an end system, but also acts as a router, that sends packets to desired nodes. The ad hoc are expected to do assignments, which the infrastructure can't do. Ad hoc networks are mostly used by military, rescue mission team, taxi driver. Their works can't rely on a infrastructure's network. As an illustrative example, imagine firefighters put out hazardous fire in a big forest. They have to communicate each other, but establishing a infrastructure or cabling in such area is impossible or too expensive.

The main problems in ad hoc networks are routing and characteristic of wireless communication. In infrastructure's networks a node can communicate with all nodes in the same cell. In ad hoc a node can communicate only with nodes in its area. This node can communicate with other nodes, but a routing algorithm is necessary. Unlike wired communication, wireless networks have transmission problem with data transmission such as, possibility of asymmetric connections and higher interferences.





2. Ad hoc networks



Fig. 2. Transmission area in ad hoc

The word ad hoc is from Latin and means "for this (only)". In the case of computer networks, the ad hoc networks mean wireless network without infrastructure, they can be called spontaneous network. One Way to understand ad hoc networks



is by comparing them with infrastructure based wireless networks, such as cellular network and WLAN. In the infrastructure based wireless networks a node can only send a packet to a destination node only via access point (in cellular network like GSM, it is called base station). The access point establishes a network area and only the nodes in this area can use access point's services.

#### A. Architecture

The wireless networks can be categorized based on their system architecture into two basically version. The one is Infrastructure (Figure 3a) and second is ad-hoc network. The biggest different of them is infrastructure networks consist of access point and nodes, meanwhile the ad hoc networks are independent from access point.



Fig. 3. Standard IEEE802.11 protocol architecture

### 1) IEEE 802.11 used for Ad Hoc Networks

The IEEE Standard 802.11 (IEEE, 1999) describes common family of wireless LANs. The standard specifies physical layer (PHY) and medium access control (MAC) of wireless transmission. The main purpose of this standard was the specification of simple and robust wireless LANs. The standard is expected to support the energy conservation of the mobile terminal, consideration of hidden terminal, and possibility of a global license-free service.

#### • PHY layer

802.11 supports 3 versions of physical layer. The one is infrared, and two others use radio transmission to transmit data (typically in 2,4 GHz). They are Frequency Hopping Spread Spectrum (FHSS) and Direct Sequence Spread Spectrum (DSSS). All of these versions provide Clear Assessment-(CCA-) signal and inform, if the medium is free. The physical layer also provides Service Access Point (SAP) with a data rate 1 or 2 Mbit/s.

#### • MAC Layer

MAC layer is responsible for many assignment. The most important assignment is obviously control the media access. This layer supports also roaming, authentication, and power management. The basic MAC layer's services are sported asynchronous data service and optional time-bounded service. The IEEE-Standard 802.11 for ad hoc provide only asynchronous data service.

## *B. Ad-hoc on-demand distance vector routing protocol* (*AODV*)

AODV is a MANET routing protocol which employs an ondemand approach to find routes, that is, a route is discovered only when it is needed by a source station. AODV uses sequence numbers to make certain freshness of routes. AODV employs route request (RREQ) packets broadcasted through the group of connections to discover the paths needed by a source station. it allows stations to find routes very fast for new destinations, and does not need stations to store routes to destinations which are not moving. AODV aids stations to operate in response to an alteration in network topology and link breakages quickly and the AODV operation is loop-free



Fig. 4. Ad-hoc on-demand distance vector routing protocol

#### 3. Wi-Fi hidden terminal and ad hoc grids

#### A. Wi-Fi Network

Wi-Fi is stands for wireless fidelity. Wi-Fi uses the 802.11 standard; it was developed by the Institute of Electrical and Electronics Engineers (IEEE) in 1997. Wi-Fi was superseded by the extensions 802.11a and 802.11b, and later by 802.11g. Wi-Fi technology uses radio communication and operating at a frequency of 2.4GHz.

- 1) Issues in Wi-Fi networks:
  - Hidden Terminal Problem
  - Exposed Terminal Problem

2) The Hidden Terminal Problem:

The hidden node/ terminal problem found at a point to multipoint network and it is defined as being one in which three or more nodes are present. Let there are three nodes: node A, node B and node C.

- A and C cannot hear each other.
- A sends to B, C cannot receive A.
- C wants to send B, C senses a free medium.
- Collision occurs at B.
- A cannot receive the collision.
- A is hidden for C.

The following figure-03 shows the Hidden Terminal Problem using node A, B and C



Fig. 5. Hidden terminal problem



### International Journal of Research in Engineering, Science and Management Volume-3, Issue-1, January-2020 www.ijresm.com | ISSN (Online): 2581-5792

### 3) Solution of hidden terminal problem

The solution of hidden terminal problem is as follows.

- When A wants to send a packet to B, A first sends a Request-to-send (RTS) to B.
- On receiving RTS, B responds by sending Clear-to-Send (CTS).
- When C overhears a CTS, it keeps quiet for the duration of the transfer.
- Transfer duration is included in both RTS and CTS.
- RTS and CTS are short frames, reduces collision chance.

# *B.* The other methods that can be employed to solve hidden terminal problem are

- Increase Transmitting Power from the Nodes.
- Use unidirectional antennas.
- Remove obstacles.
- Move the node.
- Use protocol enhancement software.
- Use antenna diversity.

#### C. Effect of hidden terminal problem

If one node hidden to another then the re-transmission will increase. It also increases the delay and decrease the throughput.

#### 4. Simulation code and results

Classical hidden terminal problem and its RTS/CTS solution. Topology: [node 0] <-- -50 dB --> [node 1] <-- -50 dB --> [node 2]

This example illustrates the use of

- Wi-Fi in ad-hoc mode
- Matrix propagation loss model
- Use of On Off Application to generate CBR stream
- IP flow monitor

Run single 10 seconds experiment with enabled or disabled RTS/CTS mechanism

### Create 3 nodes

- Place nodes somehow, this is required by every wireless simulation
- Create propagation loss matrix
- Create & setup Wi-Fi channel
- Install wireless devices
- Install TCP/IP stack & assign IP addresses
- Install applications: two CBR streams each saturating the channel
- Install Flow Monitor on all nodes
- Run simulation for 10 seconds
- Print per flow statistics
- Cleanup

#### 5. Result

Hidden station experiment with RTS/CTS disabled: Flow 1 (10.0.0.1 -> 10.0.0.2) Tx Bytes: 3441480 Tx Bytes: 3441480 Tx Bytes: 3441480 Tx Bytes: 362752 Throughput: 0.233557 Mbps Flow 2 (10.0.0.3 -> 10.0.0.2) Tx Packets: 2411 Tx Bytes: 346298 Rx Bytes: 346298 Rx Bytes: 346298 Rx Bytes: 346298 Hidden tation.periment with RTS/CTS enabled: Flow 1 (10.0.1 -> 10.0.0.2) Tx Packets: 340 FLOW 1 (10.0.1 -> 10.0.0.2) Tx Bytes: 341480 Tx Bytes: 341480 Tx Bytes: 341480 Tx Bytes: 340 FLOW 2 (10.0.0.3 -> 10.0.0.2) Tx Packets: 240 Tx Bytes: 341480 Tx Bytes: 341480 Tx Bytes: 341480 Tx Contexts: 240 Tx Bytes: 34290 Throughput: 0.69928 Mbps FLOW 2 (10.0.0.3 -> 10.0.0.2) Tx Packets: 344290 Tx Bytes: 344290 Tx Byt

1) Creating of 4 nodes

#### Hidden station experiment with RTS/CTS disabled: Flow 1 (10.0.0.1 -> 10.0.0.2) Tx Packets: 3571 Tx Bytes: 509388 TxOffered: 4.53279 Mbps Rx Pytes: 430404 Throughput: 0.545813 Mbps Flow 2 (10.0.0.3 -> 10.0.0.2) Tx Packets: 2232 Tx Packets: 183315 Mbps Rx Packets: 82 Rx Packets: 82 Rx Packets: 82 Rx Packets: 82 Troughput: 0.104085 Mbps Flow 3 (10.0.0.4 -> 10.0.0.2) Tx Bytes: 1912092 Tx Bytes: 1912092 Tx Bytes: 1927 Rx Packets: 10 Rx Packets

Hidden station exp	eriment with RTS/CTS enabled:
Flow 1 (10.0.0.1 -:	> 10.0.0.2)
Tx Packets: 3571	
Tx Bytes: 5099	388
TxOffered: 4.53	279 Mbps
Rx Packets: 559	
Rx Bytes: 7982	52
Throughput: 0.70	9557 Mbps
Flow 2 (10.0.0.3 -:	> 10.0.0.2)
Tx Packets: 2232	
Tx Bytes: 3187	296
TxOffered: 2.83	315 Mbps
Rx Packets: 247	
Rx Bytes: 3527	16
Throughput: 0.31	3525 Mbps
Flow 3 (10.0.0.4 -:	> 10.0.0.2)
Tx Packets: 1339	
Tx Bytes: 1912	092
TxOffered: 1.69	964 Mbps
Rx Packets: 126	
Rx Bytes: 1799	28
Throughput: 0 15	9936 Mbos

2) Creating of 5 nodes

Hidde	n station	a experiment with	RTS/CTS	disabled:
Flow	1 (10.0.0	0.1 -> 10.0.0.2)		
T×	Packets:	2410		
T×	Bytes:	3441480		
TXO	ffered:	3.05909 Mbps		
RX	Packets:	15		
R×	Bytes:	21420		
Thr	oughput:	0.01904 Mbps		
Flow	2 (10.0.0	$0.3 \rightarrow 10.0.0.2)$		
тх	Packets:	2411		
Тх	Bytes:	3442908		
Тхо	ffered:	3.06036 Mbps		
R×	Packets:	16		
RX	Bytes:	22848		
Thr	oughput:	0.0203093 Mbps		
Flow	3 (10.0.0	$0.4 \rightarrow 10.0.0.2$		
тх	Packets:	2411		
тх	Bytes:	3442908		
тхо	ffered:	3.06036 Mbps		
R×	Packets:	19		
RX	Bytes:	27132		
Thr	oughput:	0.0241173 Mbps		
Flow	4 (10.0.0	0.5 -> 10.0.0.2)		
T×	Packets:	2411		
т×	Bytes:	3442908		
тхо	ffered:	3.06036 Mbps		
R×	Packets:	26		
R×	Bytes:	37128		
Thr	oughput:	0.0330027 Mbps		





#### 3) Creating of 6 nodes



Graphs: No. of Nodes vs. Throughput



#### Nodes vs. throughput"

The above graph illustrates that as the number of nodes increases the throughput gradually decreases. As theoretically we know that through put (efficiency) is inversely proportional to the number of nodes present in the network. Practically observed in the above graph. As the nodes increases in the network traffic increases.

#### Data rate vs Throughput:



#### Data rate vs. throughput:

Theoretically we know that as the data rate is increases throughput decreases because the speed of the source is more than the speed of the destination the packet drops are more. As the data rate is less for the destination it cannot handle the packets which are from the source so they are dropped. By this the efficiency decreases gradually as shown in the graph.

#### 6. Conclusion

Mobility in the wireless networks is very popular nowadays. Many peoples in the street walk and are using small devices like PDA, laptops, or phone to communicate, listening a music, write SMS, exchanging data with other people near them, etc. The wireless infrastructure networks support great mobility and very popular among the folks. But this kind of networks is centralized, not flexible, and sometimes too expensive. If a infrastructure is defect, the cell established by this infrastructure will be gone too. The nodes in this cell can't communicate again. The presence of ad hoc networks covers the infrastructure's weakness. Since the ad hoc networks are independent from infrastructure, the nodes must be able to work together to establish a greater network. They have multi hop the packet, if they have to send a packet to a destination node outside their transmission range.

Therefore, routing algorithms are the main challenge. Since the nodes are mobile, link between nodes are not symmetric, and the topology are always changed, the routing algorithms used in wired network must be modified or must be invented. The ad hoc networks still have to deal with wireless problems, such as security and higher error rate. Especially WANETs have to consider their power supply, since they are not supported with fixed power supply. At least, the ad hoc networks are developed not to replace the infrastructure networks. With the great number of wireless user and frequency limitations, it is unlikely possible to control independent network. The ad hoc networks can replace the infrastructure networks only for a short time and are used for some specific situation, in which the infrastructure networks fail. A combination of both networks like figure 13 can offer great output, like connection between ad hoc and Internet.

#### References

- [1] J. Schiller, Mobilkommunikation. Addison-Wesley Verlag, 2000.
- [2] J. Schiller, "Wireless transmission," 2008.
- [3] K. U. R. Khan, R. U. Zaman, and A. V. Reddy, "Performance comparison of on-demand and table driven ad hoc routing protocols using netuns," pp. 336–341, 2008.
- [4] S. Toumpis and D. Toumpakaris, "Wireless ad hoc networks and related topologies: applications and research challenges," e & i Elektro technikund Informations technik, vol. Volume 123, pp. 232–241, 2006.
- [5] M. Günes, B. Blywis, and F. Juraschek, "Concept and design of the hybrid distributed embedded systems testbed," August 2008.
- [6] Perkins and C. E., Ad hoc networking. Addison-Wesley Verlag, 2001.
- [7] "The official bluetooth."
- [8] C. Perkins and P. Bhagwat, "Highly dynamic destination-sequenced distance-vector routing (dsdv) for mobile computers," 1994.



- [9] D. Johnson, Y. Hu, and D. Maltz, "The dynamic source routing protocol (dsr) for mobile ad hoc networks for ipv4," Network Working Group Request for Comments: 4728, 2007.
- [10] C. Perkins and E. Belding-Royer, "Ad hoc on-demand distance vector (aodv) routing," Network Working Group Request for Comments: 3561, 2003.
- [11] J. Schiller, "Network protocols / mobile ip," 2008.
- [12] "National institute of standards and technology: http://www.antd.nist.gov/index.shtml."

- [13] D. Wang, M. Hu, and H. Zhi, "A survey of secure routing in ad hoc networks," The Ninth International Conference on Web-Age Information, pp. 482–486, 2008.
- [14] T. Karygiannis and L. Owens, "Wireless network security 802.11, bluetooth and handheld devices," 2002.
- [15] S. Mueller, R. P. Tsang, and D. Ghosal, "Multipath routing in mobile ad hoc networks: Issues and challenges," Performance Tools and Applications to Networked Systems, pp. 209–234, 2004.